

15-213

"The course that gives CMU its Zip!"

Machine-Level Programming II:

Control Flow

Jan 29, 2004

Topics

- Condition Codes
 - Setting
 - Testing
- Control Flow
 - If-then-else
 - Varieties of Loops
 - Switch Statements

class06.ppt

Condition Codes

Single Bit Registers

CF	Carry Flag	SF	Sign Flag
ZF	Zero Flag	OF	Overflow Flag

Implicitly Set By Arithmetic Operations

addl Src, Dest

C analog: $t = a + b$

- CF set if carry out from most significant bit
 - Used to detect unsigned overflow
- ZF set if $t == 0$
- SF set if $t < 0$
- OF set if two's complement overflow
 - $(a>0 \&& b>0 \&& t<0) \mid\mid (a<0 \&& b<0 \&& t>=0)$

Not Set by leal instruction

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Setting Condition Codes (cont.)

Explicit Setting by Compare Instruction

cmp1 Src2,Src1

- cmp1 b,a like computing $a-b$ without setting destination
- CF set if carry out from most significant bit
 - Used for unsigned comparisons
- ZF set if $a == b$
- SF set if $(a-b) < 0$
- OF set if two's complement overflow
 - $(a>0 \&& b<0 \&& (a-b)<0) \mid\mid (a<0 \&& b>0 \&& (a-b)>0)$

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Setting Condition Codes (cont.)

Explicit Setting by Test instruction

test1 Src2,Src1

- Sets condition codes based on value of Src1 & Src2
 - Useful to have one of the operands be a mask
- test1 b,a like computing $a\&b$ without setting destination
- ZF set when $a\&b == 0$
- SF set when $a\&b < 0$

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Reading Condition Codes

SetX Instructions

- Set single byte based on combinations of condition codes

SetX	Condition	Description
sete	ZF	Equal / Zero
setne	~ZF	Not Equal / Not Zero
sets	SF	Negative
setsns	~SF	Nonnegative
setg	~(SF^OF) & ~ZF	Greater (Signed)
setge	~(SF^OF)	Greater or Equal (Signed)
setl	(SF^OF)	Less (Signed)
setle	(SF^OF) ZF	Less or Equal (Signed)
seta	~CF & ~ZF	Above (unsigned)
setb	CF	Below (unsigned)

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Reading Condition Codes (Cont.)

SetX Instructions

- Set single byte based on combinations of condition codes
- One of 8 addressable byte registers
 - Embedded within first 4 integer registers
 - Does not alter remaining 3 bytes
 - Typically use `movzbl` to finish job

```
int gt (int x, int y)
{
    return x > y;
}
```

Body

```
movl 12(%ebp),%eax # eax = y
cmpl %eax,8(%ebp) # Compare x : y
setg %al # al = x > y
movzbl %al,%eax # Zero rest of %eax
```

%eax	%ah	%al
%edx	%dh	%dl
%ecx	%ch	%cl
%ebx	%bh	%bl
%esi		
%edi		
%esp		
%ebp		

Note inverted ordering!

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Jumping

jX Instructions

- Jump to different part of code depending on condition codes

jX	Condition	Description
jmp	1	Unconditional
je	ZF	Equal / Zero
jne	~ZF	Not Equal / Not Zero
js	SF	Negative
jns	~SF	Nonnegative
jg	~(SF^OF) & ~ZF	Greater (Signed)
jge	~(SF^OF)	Greater or Equal (Signed)
jl	(SF^OF)	Less (Signed)
jle	(SF^OF) ZF	Less or Equal (Signed)
ja	~CF & ~ZF	Above (unsigned)
jb	CF	Below (unsigned)

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Conditional Branch Example

```
max:
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%edx
    movl 12(%ebp),%eax
    cmpl %eax,%edx
    jle L9
    movl %edx,%eax
L9:
    movl %ebp,%esp
    popl %ebp
    ret
```

Set Up

Body

Finish

Conditional Branch Example (Cont.)

```
int goto_max(int x, int y)
{
    int rval = y;
    int ok = (x <= y);
    if (ok) goto done;
    rval = x;
done:
    return rval;
}
```

- C allows “goto” as means of transferring control
 - Closer to machine-level programming style
- Generally considered bad coding style

```
movl 8(%ebp),%edx      # edx = x
movl 12(%ebp),%eax     # eax = y
cmpl %eax,%edx        # x : y
jle L9                  # if x <= y goto done
movl %edx,%eax         # eax = x
L9:                      # done:
```

Skipped when x ≤ y

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“Do-While” Loop Example

C Code

```
int fact_do(int x)
{
    int result = 1;
    do {
        result *= x;
        x = x-1;
    } while (x > 1);

    return result;
}
```

Goto Version

```
int fact_goto(int x)
{
    int result = 1;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
    return result;
}
```

- Use backward branch to continue looping
- Only take branch when “while” condition holds

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“Do-While” Loop Compilation

Goto Version

```
int
fact_goto(int x)
{
    int result = 1;

loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;

    return result;
}
```

Assembly

Registers
%edx x
%eax result

```
_fact_goto:
    pushl %ebp          # Setup
    movl %esp,%ebp       # Setup
    movl $1,%eax         # eax = 1
    movl 8(%ebp),%edx   # edx = x

L11:
    imull %edx,%eax    # result *= x
    decl %edx           # x--
    cmpl $1,%edx         # Compare x : 1
    jg L11               # if > goto loop

    movl %ebp,%esp        # Finish
    popl %ebp             # Finish
    ret                  # Finish
```

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General “Do-While” Translation

C Code

```
do
    Body
    while (Test);
```

Goto Version

```
loop:
    Body
    if (Test)
        goto loop;
```

- Body can be any C statement
 - Typically compound statement:

```
{
    Statement1;
    Statement2;
    ...
    Statementn;
}
```

- Test is expression returning integer
= 0 interpreted as false ≠ 0 interpreted as true

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“While” Loop Example #1

C Code

```
int fact_while(int x)
{
    int result = 1;
    while (x > 1) {

        result *= x;
        x = x-1;
    };

    return result;
}
```

- Is this code equivalent to the do-while version?
- Must jump out of loop if test fails

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First Goto Version

```
int fact_while_goto(int x)
{
    int result = 1;
    loop:
        if (!(x > 1))
            goto done;
        result *= x;
        x = x-1;
        goto loop;
    done:
        return result;
}
```

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Actual “While” Loop Translation

C Code

```
int fact_while(int x)
{
    int result = 1;
    while (x > 1) {
        result *= x;
        x = x-1;
    };
    return result;
}
```

- Uses same inner loop as do-while version
- Guards loop entry with extra test

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Second Goto Version

```
int fact_while_goto2(int x)
{
    int result = 1;
    if (!(x > 1))
        goto done;
    loop:
        result *= x;
        x = x-1;
        if (x > 1)
            goto loop;
    done:
        return result;
}
```

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General “While” Translation

C Code

```
while (Test)
    Body
```



Do-While Version

```
if (!Test)
    goto done;
do
    Body
    while (Test);
done:
```



Goto Version

```
if (!Test)
    goto done;
loop:
    Body
    if (Test)
        goto loop;
done:
```

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“For” Loop Example

```
/* Compute x raised to nonnegative power p */
int ipwr_for(int x, unsigned p)
{
    int result;
    for (result = 1; p != 0; p = p>>1) {
        if (p & 0x1)
            result *= x;
        x = x*x;
    }
    return result;
}
```

Algorithm

- Exploit property that $p = p_0 + 2p_1 + 4p_2 + \dots + 2^{n-1}p_{n-1}$
- Gives: $x^p = z_0 \cdot z_1^2 \cdot z_2^4 \cdot \dots \cdot ((z_{n-1})^2)^{2^{n-1}}$
- $z_i = 1$ when $p_i = 0$
- $z_i = x$ when $p_i = 1$
- Complexity O(log p)

Example

$$3^{10} = 3^2 \cdot 3^8$$

$$= 3^2 \cdot ((3^2)^2)^2$$

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ipwr Computation

```
/* Compute x raised to nonnegative power p */
int
ipwr_for(int x, unsigned p)
{
    int result;
    for (result = 1; p != 0; p = p>>1) {
        if (p & 0x1)
            result *= x;
        x = x*x;
    }
    return result;
}
```

result	x	p
1	3	10
1	9	5
9	81	2
9	6561	1
531441	43046721	0

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“For” Loop Example

```
int result;
for (result = 1;
     p != 0;
     p = p>>1)
{
    if (p & 0x1)
        result *= x;
    x = x*x;
}
```

General Form
for (Init; Test; Update)
Body

Init
result = 1

Test
p != 0

Update
p = p >> 1

Body
`{
 if (p & 0x1)
 result *= x;
 x = x*x;
}`

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“For”→“While”

For Version

```
for (Init; Test; Update)
    Body
```

While Version

```
Init;
while (Test) {
    Body
    Update ;
}
```

Do-While Version

```
Init;
if (!Test)
    goto done;
do {
    Body
    Update ;
} while (Test)
done:
```

Goto Version

```
Init;
if (!Test)
    goto done;
loop:
    Body
    Update ;
    if (Test)
        goto loop;
done:
```

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“For” Loop Compilation

Goto Version

```
Init;
if (!Test)
    goto done;
loop:
    Body
    Update ;
    if (Test)
        goto loop;
done:
```

```
result = 1;
if (p == 0)
    goto done;
loop:
    if (p & 0x1)
        result *= x;
    x = x*x;
    p = p >> 1;
    if (p != 0)
        goto loop;
done:
```

Init

result = 1

Test

p != 0

Update

p = p >> 1

Body

```
{  
    if (p & 0x1)  
        result *= x;  
    x = x*x;  
}
```

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```

typedef enum
{ADD, MULT, MINUS, DIV, MOD, BAD}
    op_type;

char unparse_symbol(op_type op)
{
    switch (op) {
    case ADD :
        return '+';
    case MULT:
        return '*';
    case MINUS:
        return '-';
    case DIV:
        return '/';
    case MOD:
        return '%';
    case BAD:
        return '?';
    }
}

```

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Switch Statements

Implementation Options

- Series of conditionals
 - Good if few cases
 - Slow if many
- Jump Table
 - Lookup branch target
 - Avoids conditionals
 - Possible when cases are small integer constants
- GCC
 - Picks one based on case structure
- Bug in example code
 - No default given

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Jump Table Structure

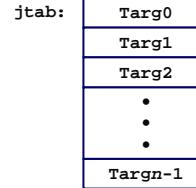
Switch Form

```

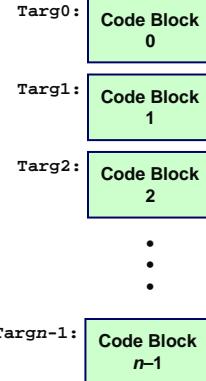
switch(op) {
    case val_0:
        Block 0
    case val_1:
        Block 1
    ...
    case val_n-1:
        Block n-1
}

```

Jump Table



Jump Targets



Approx. Translation

```

target = JTab[op];
goto *target;

```

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Switch Statement Example

Branching Possibilities

```

typedef enum
{ADD, MULT, MINUS, DIV, MOD, BAD}
    op_type;

char unparse_symbol(op_type op)
{
    switch (op) {
        ...
    }
}

```

Setup:

```

unparse_symbol:
    pushl %ebp      # Setup
    movl %esp,%ebp  # Setup
    movl 8(%ebp),%eax # eax = op
    cmpl $5,%eax   # Compare op : 5
    ja .L49         # If > goto done
    jmp *.L57(%eax,4) # goto Table[op]

```

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Enumerated Values

ADD	0
MULT	1
MINUS	2
DIV	3
MOD	4
BAD	5

Assembly Setup Explanation

Symbolic Labels

- Labels of form .LXX translated into addresses by assembler

Table Structure

- Each target requires 4 bytes
- Base address at .L57

Jumping

- jmp .L49
- Jump target is denoted by label .L49
- jmp *.L57(%eax,4)
- Start of jump table denoted by label .L57
- Register %eax holds op
- Must scale by factor of 4 to get offset into table
- Fetch target from effective Address .L57 + op*4

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Jump Table

Table Contents

```
.section .rodata
.align 4
.L57:
.long .L51 #Op = 0
.long .L52 #Op = 1
.long .L53 #Op = 2
.long .L54 #Op = 3
.long .L55 #Op = 4
.long .L56 #Op = 5
```

Enumerated Values

ADD	0
MULT	1
MINUS	2
DIV	3
MOD	4
BAD	5

Targets & Completion

```
.L51:
    movl $43,%eax # '+'
    jmp .L49
.L52:
    movl $42,%eax # '**'
    jmp .L49
.L53:
    movl $45,%eax # '-'
    jmp .L49
.L54:
    movl $47,%eax # '//'
    jmp .L49
.L55:
    movl $37,%eax # '%'
    jmp .L49
.L56:
    movl $63,%eax # '?'
    # Fall Through to .L49
```

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Switch Statement Completion

```
.L49:          # Done:
    movl %ebp,%esp      # Finish
    popl %ebp      # Finish
    ret             # Finish
```

Puzzle

- What value returned when op is invalid?

Answer

- Register %eax set to op at beginning of procedure
- This becomes the returned value

Advantage of Jump Table

- Can do k-way branch in O(1) operations

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Object Code

Setup

- Label .L49 becomes address 0x804875c
- Label .L57 becomes address 0x8048bc0

```
08048718 <unparse_symbol>:
8048718: 55          pushl  %ebp
8048719: e9 45 08    movl   %esp,%ebp
804871b: 8b 45 08    movl   0x8(%ebp),%eax
804871e: f8 05       cmpl   $0x5,%eax
8048721: 77 39       ja     804875c <unparse_symbol+0x44>
8048723: ff 24 85 c0 8b jmp    *0x8048bc0(%eax,4)
```

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Object Code (cont.)

Jump Table

- Doesn't show up in disassembled code
 - Can inspect using GDB
- gdb code-examples
- ```
(gdb) x/6xw 0x8048bc0
● Examine 6 hexadecimal format "words" (4-bytes each)
● Use command "help x" to get format documentation
```

0x8048bc0 <\_fini+32>:

```
0x08048730
0x08048737
0x08048740
0x08048747
0x08048750
0x08048757
```

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## Extracting Jump Table from Binary

### Jump Table Stored in Read Only Data Segment (.rodata)

- Various fixed values needed by your code

### Can examine with objdump

```
objdump code-examples -s --section=.rodata
▪ Show everything in indicated segment.
```

### Hard to read

- Jump table entries shown with reversed byte ordering

Contents of section .rodata:

```
8048bc0 30870408 37870408 40870408 47870408 0...7...@...G...
8048bd0 50870408 57870408 46616374 28256429 P...W...Fact(%d)
8048be0 203d2025 6c640a00 43686172 203d2025 = %ld..Char = %
...
```

- E.g., 30870408 really means 0x08048730

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## Disassembled Targets

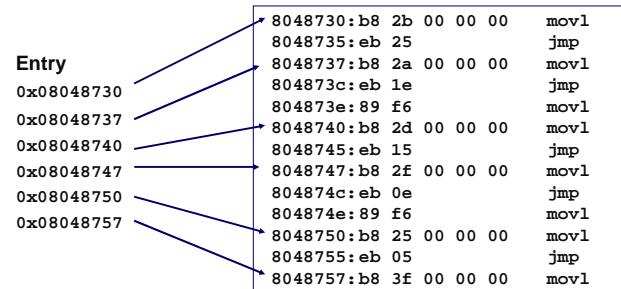
```
8048730:b8 2b 00 00 00 movl $0x2b,%eax
8048735:eb 25 jmp 804875c <unparse_symbol+0x44>
8048737:b8 2a 00 00 00 movl $0x2a,%eax
804873c:eb 1e jmp 804875c <unparse_symbol+0x44>
804873e:89 f6 movl %esi,%esi
8048740:b8 2d 00 00 00 movl $0x2d,%eax
8048745:eb 15 jmp 804875c <unparse_symbol+0x44>
8048747:b8 2f 00 00 00 movl $0x2f,%eax
804874c:eb 0e jmp 804875c <unparse_symbol+0x44>
804874e:89 f6 movl %esi,%esi
8048750:b8 25 00 00 00 movl $0x25,%eax
8048755:eb 05 jmp 804875c <unparse_symbol+0x44>
8048757:b8 3f 00 00 00 movl $0x3f,%eax
```

- movl %esi,%esi does nothing
- Inserted to align instructions for better cache performance

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## Matching Disassembled Targets



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## Sparse Switch Example

```
/* Return x/111 if x is multiple
 && <= 999. -1 otherwise */
int div111(int x)
{
 switch(x) {
 case 0: return 0;
 case 111: return 1;
 case 222: return 2;
 case 333: return 3;
 case 444: return 4;
 case 555: return 5;
 case 666: return 6;
 case 777: return 7;
 case 888: return 8;
 case 999: return 9;
 default: return -1;
 }
}
```

- Not practical to use jump table
  - Would require 1000 entries
- Obvious translation into if-then-else would have max. of 9 tests

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## Sparse Switch Code

```

movl 8(%ebp),%eax # get x
cmpb $444,%eax # x:444
je L8
jg L16
cmpl $111,%eax # x:111
je L5
jg L17
testb %eax,%eax # x:0
je L4
jmp L14
 . . .

```

- Compares x to possible case values
- Jumps different places depending on outcomes

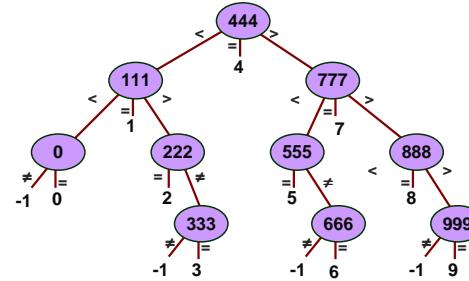
```

 . . .
L5:
 movl $1,%eax
 jmp L19
L6:
 movl $2,%eax
 jmp L19
L7:
 movl $3,%eax
 jmp L19
L8:
 movl $4,%eax
 jmp L19
 . . .

```

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## Sparse Switch Code Structure



- Organizes cases as binary tree
- Logarithmic performance

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## Summarizing

### C Control

- if-then-else
- do-while
- while
- switch

### Assembler Control

- jump
- Conditional jump

### Compiler

- Must generate assembly code to implement more complex control

### Standard Techniques

- All loops converted to do-while form
- Large switch statements use jump tables

### Conditions in CISC

- CISC machines generally have condition code registers

### Conditions in RISC

- Use general registers to store condition information
- Special comparison instructions
- E.g., on Alpha:
  - cmple \$16,1,\$1
  - Sets register \$1 to 1 when Register \$16 <= 1

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