Lecture 7

Loop Invariant Computation and Code Motion

I Finding loops

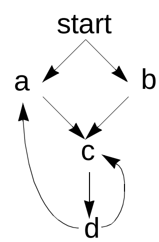
II Loop-invariant computation

III Algorithm for code motion

Reference: Muchnick 13.2

I. What is a Loop?

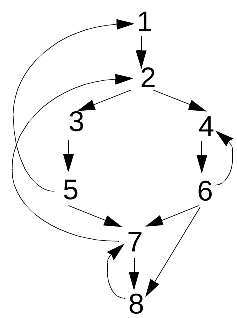
- **Goal:**
  - Define a loop in graph-theoretic terms (control flow graph)
  - Not sensitive to input syntax,
    a uniform treatment for all loops: DO, while, goto’s
- **Not every cycle is a “loop” from the optimization perspective**



- **Intuitive properties of a loop**
  - single entry point
  - edges must form at least a cycle

Formal Definitions

- **Dominators**
  - Node *d* dominates node *n* in a graph (*d dom n*) if every path from the start node to *n* goes through *d*



  - Dominators can be organized as a tree
    - a->b in the dominator tree iff a immediately dominates b

Natural Loops

- **Definitions**
  - Single entry-point: ***header***
    - a header dominates all nodes in the loop
  - A ***back edge*** is an arc whose head dominates its tail (tail -> head)
    - a back edge must be a part of at least one loop
  - The ***natural loop of a back edge*** is
    the smallest set of nodes that
    includes the head and tail of the back edge, and
    has no predecessors outside the set,
    except for the predecessors of the header.

Algorithm to Find Natural Loops

1. Find the dominator relations in a flow graph

2. Identify the back edges

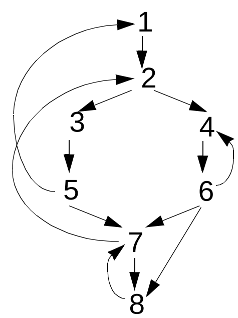
3. Find the natural loop associated with the back edge

1. Finding Dominators

- **Definition**
  - Node d dominates node n in a graph (d *dom* n) if every path from the start node to n goes through d
- **Formulated as MOP problem**
  - node d lies on all possible paths reaching node n => d *dom* n
    - Direction:
    - Values:
    - Meet operator:
    - Top:
    - Bottom:
    - Boundary condition: start/entry node =
    - Initialization for internal nodes
    - Finite descending chain?
    - Transfer function:
- **Speed:**
    - With reverse postorder, most flow graphs (reducible flow graphs) converge in 1 pass

2. Finding Back Edges

- **Depth-first spanning tree**
  - Edges traversed in a depth-first search of the flow graph form a depth-first spanning tree



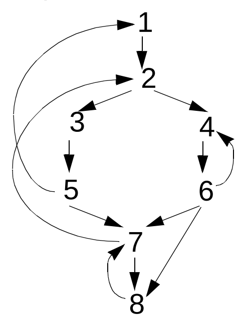
- **Categorizing edges in graph**
  - Advancing edges: from ancestor to proper descendant
  - Cross edges: from right to left
  - Retreating edges: from descendant to ancestor (not necessarily proper)

Back Edges

- **Definition**
  - Back edge: t->h, h dominates t
- **Relationships between graph edges and back edges**
- **Algorithm**
  - Perform a depth first search
  - For each retreating edge t->h, check if h is in t's dominator list
- **Most programs (all structured code, and most GOTO programs) -- reducible flow graphs**
  - retreating edges = back edges

3. Constructing Natural Loops

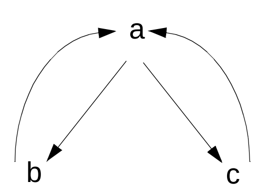
- The ***natural loop of a back edge*** is
  the smallest set of nodes that
  includes the head and tail of the back edge, and
  has no predecessors outside the set,
  except for the predecessors of the header.
- **Algorithm**
  - delete h from the flow graph
  - find those nodes that can reach t
    (those nodes plus h form the natural loop of t ->h)



Inner Loops

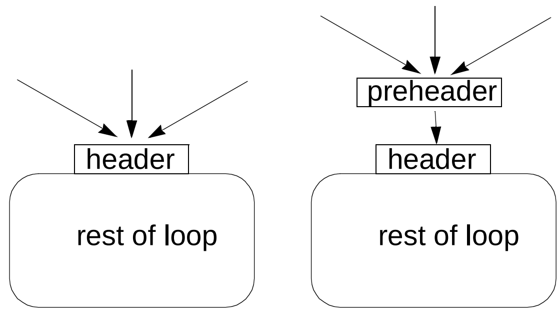
- **If two loops do not have the same header**
  - they are either disjoint, or
  - one is entirely contained (nested within) the other
    -- inner loop, one that contains no other loop.
- **If two loops share the same header**
  - Hard to tell which is the inner loop

Combine as one



Preheader

- **Optimizations often require code to be executed once before the loop**
- **Create a preheader basic block for every loop**

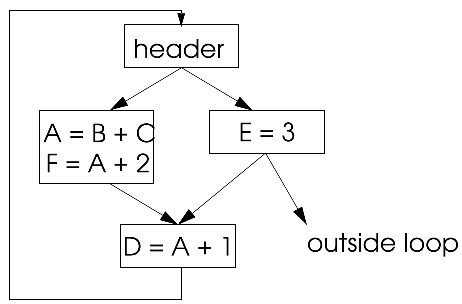


Finding Loops: Summary

- **Define loops in graph theoretic terms**
- **Definitions and algorithms for**
  - Dominators
  - Back edges
  - Natural loops

II. Loop-Invariant Computation and Code Motion

- **A loop-invariant computation:**
  - a computation whose value does not change as long as control stays within the loop
- **Code motion:**
  - to move a statement within a loop to the preheader of the loop



Algorithm

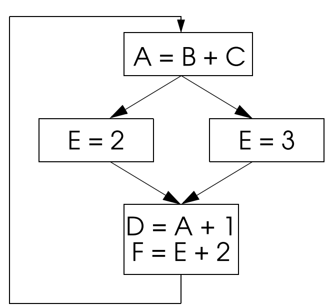
- **Observations**
  - Loop invariant
    - operands are defined outside loop or invariant themselves
  - Code motion
    - not all loop invariant instructions can be moved to preheader
- **Algorithm**
  - Find invariant expressions
  - Conditions for code motion
  - Code transformation

Detecting Loop Invariant Computation

- Compute reaching definitions
- Mark INVARIANT if
  all the definitions of B and C that reach a statement A=B+C are outside the loop
  - constant B, C?
- Repeat: Mark INVARIANT if
  - all reaching definitions of B are outside the loop, or
  - there is exactly one reaching definition for B, and it is from a loop-invariant statement inside the loop
  - similarly for C

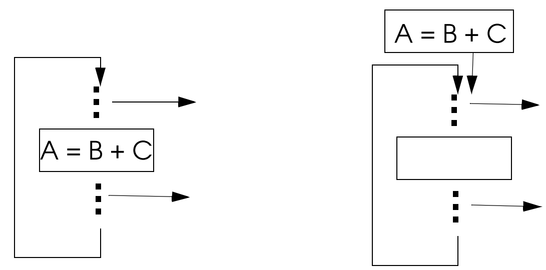
  until no changes to set of loop-invariant statements occur.

Example



III. Conditions for Code Motion

- **Correctness: Movement does not change semantics of program**
- **Performance: Code is not slowed down**



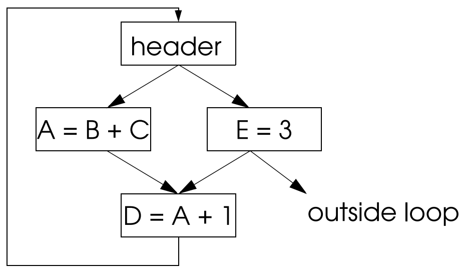
- **Basic idea: defines <u>once</u> and <u>for all</u>**
  - control flow:
  - other definitions:
  - other uses:

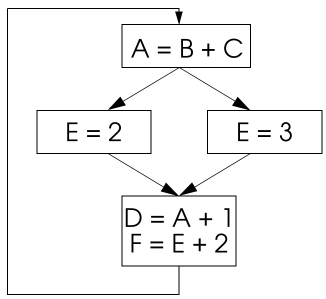
Code Motion Algorithm

Given: a set of nodes in a loop

- **Compute reaching definitions**
- **Compute loop invariant computation**
- **Compute dominators**
- **Find the exits of the loop, nodes with successor outside loop**
- **Candidate statement for code motion:**
  - loop invariant
  - in blocks that dominate all the exits of the loop
  - assign to variable not assigned to elsewhere in the loop
  - in blocks that dominate all blocks in the loop that use the variable assigned
- **Perform a depth-first search of the blocks**
  - Move candidate to preheader
    if all the invariant operations it depends on have been moved

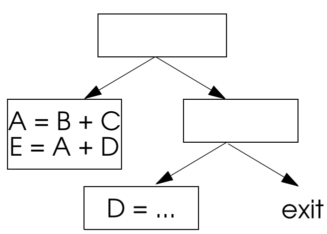
Examples





More Aggressive Optimizations

- **Gamble on: most loops get executed**
  - Can we relax constraint of dominating all exits?



- **Landing pads**

```
While p do s      ==>      if p {
                               preheader
                               repeat
                                   s
                               until not p;
                           }
```

Conclusions

- Precise definition and algorithm for loop invariant computation
- Precise algorithm for code motion
- Use of reaching definitions and dominators in optimizations