1 A few notes on Lecture 14

1.1 The distance bound

Recall that we want to bound

$$|Z_i - Z_{i-1}| = |E[f(X_1, \dots, X_{i-1}, X_i, X_{i+1}, \dots, X_n) - f(X_1, \dots, X_{i-1}, \widehat{X}_i, X_{i+1}, \dots, X_n) | X_1, X_2, \dots, X_i]|$$
(1)

Note that we do not want to just use the $2\sqrt{2}$ -Lipschitz property, since that will be too weak, and will only give us

$$\Pr[|f - Ef| \le \lambda] \le \exp\{-\lambda^2 / O(n)\}.$$

We want something much better!

Claim 1.1

$$f(X_1, \dots, X_{i-1}, X_i, X_{i+1}, \dots, X_n) - f(X_1, \dots, X_{i-1}, \hat{X}_i, X_{i+1}, \dots, X_n)$$

$$\leq 2 \big(\min_{j \neq i} d(X_i, X_j) + \min_{j \neq i} d(\hat{X}_i, X_j) \big).$$

Proof: Let $A = X_1, \ldots, X_{i-1}, X_{i+1}, \ldots, X_n$ be all the points except X_i and \hat{X}_i , and let T(A) be the optimal TSP tour on A. Note that f(A) = length(T(A)). For any point x and set S, define $d(x, S) = \min_{y \in S} d(x, y)$.

Note that if take T(A), and to it we add two edges from X_i to its closest point in A, and from \hat{X}_i to its closest point in A, then we have an Eulerian graph on the n + 1 points $A \cup \{X_i, \hat{X}_i\}$ of total length at most

$$f(A) + 2(d(X_i, A) + d(\widehat{X}_i, A)).$$

$$\tag{2}$$

Using the triangle inequality to shortcut repeated vertices gives us TSP tour of length at most (2), and hence the length of the optimal tour on $A \cup \{X_i, \hat{X}_i\}$ has length

$$f(A \cup \{X_i, X_i\}) \le f(A) + 2(d(X_i, A) + d(X_i, A)).$$
(3)

Finally, using the fact that

$$f(A) \le f(A \cup \{X_i\}) \le f(A \cup \{X_i, \hat{X}_i\})$$

$$f(A) \le f(A \cup \{\hat{X}_i\}) \le f(A \cup \{X_i, \hat{X}_i\})$$

implies that

$$|f(A \cup \{X_i\}) - f(A \cup \{\widehat{X}_i\})| \le f(A \cup \{X_i, \widehat{X}_i\}) - f(A) \le 2(d(X_i, A) + d(\widehat{X}_i, A)),$$
(4)

the last inequality using (3). This is just a rephrasing of the claim that we want to prove. **Corollary 1.2** Define $B = \{X_j \mid j > i\}$. Then

$$f(X_1, \dots, X_{i-1}, X_i, X_{i+1}, \dots, X_n) - f(X_1, \dots, X_{i-1}, \widehat{X}_i, X_{i+1}, \dots, X_n) \le 2(d(X_i, B) + d(\widehat{X}_i, B)).$$
(5)

Proof: The quantity on the right of (5) is larger than the quantity on the right of (4), since $B \subseteq A$.

Lemma 1.3

$$|Z_i - Z_{i-1}| \le 2 \left(E[d(X_i, B) \mid X_i] + E[d(\widehat{X}_i, B)] \right).$$
(6)

Proof: Plug in the result of Claim 1.1 into (1), and note that $d(X_i, B)$ is independent of $X_1, X_2, \ldots, X_{i-1}$, whereas $d(\hat{X}_i, B)$ is independent of all X_1, \ldots, X_i . Simplifying gives us the lemma.

1.2 The Rest of the Argument

Suppose we throw down n-i points randomly in \mathcal{U} , and define the random variable $Q_i(x)$ to be the distance of x to the closest point amongst these n-i. Let R be the set of random points, and hence $Q_i(x) = d(x, R)$. We proved that

Claim 1.4 For any $x \in \mathcal{U}$,

$$E[Q_i(x)] \le \frac{O(1)}{\sqrt{n-i}}.$$
(7)

Proof: This was the geometric proof, and I am going to omit it.

Hence we can upper bound both $E[d(X_i, B) \mid X_i]$ and $E[d(\widehat{X}_i, B)]$ by $\frac{O(1)}{\sqrt{n-i}}$. Finally, using (6), we get

$$|Z_i - Z_{i-1}| \le 2\left(\frac{O(1)}{\sqrt{n-i}} + \frac{O(1)}{\sqrt{n-i}}\right)$$
(8)

This implies that we can set $c_i = \frac{O(1)}{\sqrt{n-i}}$ in Azuma's inequality, which is much better than the bound that we get just plugging in the $2\sqrt{2}$ -Lipschitz-ness of f. Now $\sum_i c_i^2 = O(\log n)$, and hence we get

$$\Pr[|f - Ef| \le \lambda] \le \exp\{-\lambda^2 / O(\log n)\},\tag{9}$$

as claimed.